



TITLE:

On Forming a Series-Parallel Graph by Removing Nodes of a Planar Graph (Studies on Computational Complexities and Related Topics)

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CITATION:

WATANABE, TOSHIMASA ...[et al.]. On Forming a Series-Parallel Graph by Removing Nodes of a Planar Graph (Studies on Computational Complexities and Related Topics). 数理解析研究所講究録 1980, 381: 22-41

ISSUE DATE:

1980-04

URL:

<http://hdl.handle.net/2433/104822>

RIGHT:

On Forming a Series-parallel Graph by Removing Nodes of a Planar Graph

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1. Introduction.

Let P_{SPNR} denote the problem of determining, for a graph $G=(V,E)$ and a nonnegative integer k , whether or not a series-parallel graph can be formed by removing k or fewer nodes of G , where V and E are the sets of nodes and of edges of G , respectively. The authors have already proved in [15] that P_{SPNR} is NP-complete. In this paper, we show that P_{SPNR} remains NP-complete even when the domain is restricted to planar graphs. From now on, let P_{PSPNR} denote P_{SPNR} with the constraint that G is planar.

Technical terms and notations not specified in this paper can be identified in [2],[7],[8] and [9].

In stead of a thorough description of the formal requirements of a proof of NP-completeness, we describe the two steps (i) and (ii) required in proving that a particular problem P_X is NP-complete.

(i) Prove that $P_X \in \mathcal{NP}$, the class of problems that can be solved in polynomial time by a nondeterministic Turing machine.

(ii) Prove that some known NP-complete problem P_X , can be polynomially transformed into P_X ($P_X \propto P_X$ for short), in such a way that any polynomial-time algorithm for solving P_X could be used to solve P_X , in polynomial time.

We will omit verification of (i) from our proofs since there

exists an algorithm that determines in polynomial time whether or not a given graph is series-parallel (For example, see [3] or [12]). Thus our proofs will focus on the transformation required by (ii).

2. Preliminaries.

Let $G=(V,E)$ and $X \subset V$. Let us denote

$$E(X) = \{ (u,v) \in E : \{u,v\} \cap X \neq \emptyset \}$$

for $X = \{v_{i_1}, \dots, v_{i_x}\}$ ($x = |X|$). We call

$$G-X = (V-X, E-E(X))$$

the graph formed by removing nodes v_{i_1}, \dots, v_{i_x} (or removing X) of G . Let us denote

$$\delta(G) = \max \{ \delta_G(v) : v \in V \},$$

where $\delta_G(v)$ denotes the node degree of v in G . Let

$$P(u,v) = (V_{P(u,v)}, E_{P(u,v)})$$

denote a path of length $|E_{P(u,v)}| \geq 2$, a subgraph of the graph G in consideration, connecting two nodes u and v of G . $P(u,v)$ is said to be a disjoint path if and only if it satisfies the condition that no node in $V_{P(u,v)} - \{u,v\}$ is (and shall be) contained in any other path except $P(u,v)$ (That is, any node of $P(u,v)$ except u and v is of node degree 2 in any graph in consideration).

For two nodes u and v , let

$$L(u,v) = (V_{L(u,v)}, E_{L(u,v)})$$

denote the graph which consists of $(\alpha+1)$ disjoint paths of length 2 each of which connects u and v , where

$$V_{L(u,v)} = \{ b_i(u,v) : i=0, \dots, \alpha \} \cup \{u,v\}$$

and

$$E_{L(u,v)} = \{ (u, b_i(u,v)), (v, b_i(u,v)) : i=0, \dots, \alpha \}.$$

$L(u,v)$ is said to be a band of width α connecting u and v .

Remark 1.

Let $L(u,v)$ be a band of width α connecting u and v . Then there exists at least one disjoint path even if α or fewer nodes of $L(u,v)$ except u and v are removed. If we assume $\alpha \geq 2$, then $L(u,v) - \{u\}$ ($L(u,v) - \{v\}$) contains at least three blocks (maximal nonseparable components) each of which consists of a single edge having v (u) as a cutpoint.

A thinning is an operation that deletes one of two multiple edges e_1 and e_2 . A shrinking is an operation that contracts one of two edges (u,v) and (v,w) satisfying that v is of node degree 2. A reduction is to repeat a thinning and/or a shrinking a number of times. A graph G is said to be series-parallel (s-p for short) if and only if there exists a reduction that leads G to a single edge. A well-known characterization of s-p graphs has been given by R. J. Duffin in [3].

Theorem 1 [3].

Let G be nonseparable. Then G is s-p if and only if G has no subgraph from which K_4 can be formed by contracting and/or deleting a number of edges.

If G has such a subgraph as mentioned in Theorem 1 then we say that G has a subgraph reducible to K_4 ($G \supset K_4$ for short).

We now turn our attention to the node cover problem. A node cover for $G=(V,E)$ is a subset S of V such that any edge of G is incident upon some node of S . Given a graph $G=(V,E)$ and a non-negative integer k , let P_{NC} , called the node cover problem, denote the problem of determining whether or not G has a node cover

s of size $|S| \leq k$ (We can assume that G is a simple graph).
 It is well-known that P_{NC} is NP-complete (For example, see [9]).

In this paper we state all lemmas without proofs which are given in [16].

Lemma 1.

P_{C-QPNC} , the node cover problem for simple planar connected cubic graphs, is NP-complete.

3. Forming a series-parallel graph by removing nodes of a planar graph.

Throughout this section let k be nonnegative integer and $G=(V,E)$ be a cubic planar connected graph, where

$$V = \{v_1, \dots, v_n\} \quad (n = |V|) \text{ and } E = \{e_1, \dots, e_m\} \quad (m = |E|).$$

Represent G on a plane and fix this representation. Let

$$F_G = \{f_0, f_1, \dots, f_{r-1}\} \quad (f_0 \text{ is the infinite face of } G)$$

denote the set of all faces of G . We can assume that

$$n \geq 4, m \geq 6 \text{ and } r \geq 4.$$

Beginning with the fixed representation of G we construct the planar graphs $G_i=(V_i, E_i)$ ($i=1,2$) and $G_1'=(V_1', E_1')$ by the following procedures (1), (2), (3) and (4).

(1) Construct the geometric dual $G^*=(V^*, E^*)$ of G , and then determine a spanning tree $T^*=(V_{T^*}, E_{T^*})$ of G^* . There exists one to one correspondence between E and E^* . We denote this correspondence by $e^* \in E^*$ for $e \in E$. Let

$$E_T = \{e \in E : e^* \in E_{T^*}\}.$$

Then $|E_T| = r-1$.

(2) Let G_1 be the graph obtained by replacing each edge $e=(u,v) \in E_T$ with a disjoint path $P(u,v)=(V_{P(u,v)}, E_{P(u,v)})$ of length 4. $P(u,v)$ is called the T -path for e . Let us denote

$$V_{P(u,v)} = \{v_0=u, v_1, v_2, v_3, v_4=v\}, \text{ and}$$

$$E_{P(u,v)} = \{(v_i, v_{i+1}) : i=0, \dots, 3\}$$

(hereafter $V_{P(u,v)}$ and $E_{P(u,v)}$ are used in this sense). Let

$$A_1 = \bigcup_{(u,v) \in E_T} V_{P_a(u,v)},$$

where $V_{P_a(u,v)} = V_{P(u,v)} - \{u, v\}$.

Then $V_1 = V \oplus A_1$, where \oplus denotes the disjoint union. Let F_{G_1} denote the set of all faces of G_1 . Then there exists one to one correspondence between F_{G_1} and F_G , and we denote this correspondence by $f^{(1)} \in F_{G_1}$ for $f \in F_G$.

(3) For each edge $e=(u,v) \in E$, execute the following procedures (i), (ii) and (iii). Let f_i and $f_j \in F_G$ be the two faces (not necessarily distinct) that e touches. If e is a bridge of G then e^* is a loop of G^* . Since $e^* \notin E_{T^*}$, we have $e \notin E_T$. That is, if $e \in E_T$ then e is not a bridge of G .

(i) The case when $e \in E - E_T$.

(i-a) If e is not a bridge of G (then $f_i \neq f_j$), place two bands $L(u^i(e), v^i(e))$ and $L(u^j(e), v^j(e))$ of width $(k+r+2)$ in $f_i^{(1)}$ and $f_j^{(1)}$, respectively.

(i-b) If e is a bridge of G (then $f_i = f_j$), place two bands $L(u_a^i(e), v_a^i(e))$ and $L(u_b^i(e), v_b^i(e))$ of width $(k+r+2)$ in $f_i^{(1)}$, where we assume that one of them is located across e from the other.

(ii) The case when $e \in E_T$ (then $f_i = f_j$).

Let $V_{P(u,v)}$ be as before. Then place four bands

$L(v_t^i(e), v_{t+1}^i(e))$ ($i=0,1,2,3$) of width $(k+r+2)$ in $f_i^{(1)}$,

and similarly four bands $L(v_t^j(e), v_{t+1}^j(e))$ ($j=0,1,2,3$)

of width $(k+r+2)$ in $f_j^{(1)}$, where $v_0(e)=u^\alpha(e)$ and $v_4(e)=v^\alpha(e)$ for $\alpha=i,j$.

- (iii) For each $u \in V$ and each pair of edges $e=(u,v)$ and $g=(u,w)$ ($e, g \in E$ and $v \neq w$), both of which are incident upon u , connect $u^i(e)$ and $u^i(g)$ in terms of a band of width $(k+r+2)$ if and only if both e and g touch a face $f_i^{(1)} \in F_{G_1}$. If e (g) is a bridge of G , then we assume that $u^i(e)$ is set to either $u_a^i(e)$ or $u_b^i(e)$ (either $u_a^i(g)$ or $u_b^i(g)$).

Let G_1' denote the resulting graph. Let V_{C_i} denote the set of nodes placed in $f_i^{(1)} \in F_{G_1}$ in (i), (ii) and (iii) of (2), and let $C_i=(V_{C_i}, E_{C_i})$ denote the subgraph of G_1' induced by V_{C_i} . C_i can be changed into a single circuit if we replace each band of C_i with a disjoint path of length 2. C_i is said to be the inner ring $IR(f_i)$ of f_i (or of $f_i^{(1)}$).

(4) Construct a maximum matching $M \subset E$ of G (It is well-known that there exists a polynomial-time algorithm to obtain a maximum matching of a planar graph G (see [4])). And execute the following procedures (i), (ii) and (iii) with respect to M .

- (i) For each edge $e=(u,v) \in E_T$, let f_i and f_j ($f_i \neq f_j$) be two faces of G that e touches, where we can assume that one of them is a finite face. Then, for each $v_t \in V_{P_a}(u,v)$ ($t=1,2,3$), connect two pairs of nodes $\{v_t, v_t^i(e)\}$ and $\{v_t, v_t^j(e)\}$ by edges $(v_t, v_t^i(e))$ and $(v_t, v_t^j(e))$, respectively. A path of length 2 defined by each pair of these edges is called the c -path. Then we

say that $IR(f_i)$ and $IR(f_j)$ are T-connected. e is said to be the c-edge of these two inner rings, and then these two inner rings are said to be the c-supporters of e .

- (ii) For each edge $e=(u,v) \in M$, choose arbitrarily a face of G that e touches. Let, say, $f_i \in F_G$ denote this face. Let $g=(u,x)$ and $h=(v,y) \in E-M$ ($\{x,y\} \cap \{u,v\} = \emptyset$) denote the two edges touching f_i . Then connect four pairs of nodes $\{u, u^i(e)\}$, $\{u, u^i(g)\}$, $\{v, v^i(e)\}$ and $\{v, v^i(h)\}$ by edges $(u, u^i(e))$, $(u, u^i(g))$, $(v, v^i(e))$ and $(v, v^i(h))$, respectively. e is said to be the b-edge of $IR(f_i)$, and $IR(f_i)$ to be the supporter of e .
- (iii) For each $v \in V-M(V)$ ($M(V) = \{u, v \in V: (u,v) \in M\}$), choose arbitrarily an edge of E incident upon v and choose again arbitrarily a face of G that e touches. Let, say, $f_i \in F_G$ denote this face. Let $g=(v,x)$ and $h=(v,y) \in E$ ($x \neq y$) denote the two edges touching f_i , both of which are incident upon v . Then connect two pairs of nodes $\{v, v^i(g)\}$ and $\{v, v^i(h)\}$ by edges $(v, v^i(g))$ and $(v, v^i(h))$, respectively. v is said to be the b-node of $IR(f_i)$, and $IR(f_i)$ be the supporter of v .

Let G_2 denote the graph constructed by applying the procedure (4) to G_1' . Since it is well-known that G^* , T^* and M can be obtained in polynomial time, G_2 can be constructed in polynomial time when G is given as input. Note that G_2 is a non-separable planar graph. Figure 1 shows an example of the graph transformation from G to G_2 .

Let v be the node mentioned in (iii) of (4). Each $v \in V$ is connected to exactly one inner ring C_i in terms of a pair of edges (v, v') and (v, v'') ($v', v'' \in V_{C_i}$) both of which are intro-

duced in either (ii) or (iii) of (4). Hence $\delta_{G_2}(v)=5$. Let $\Delta(v)$ denote the subgraph of G_2 called the triangle of C_i (or of v), where it is defined by nodes v, v' and v'' , edges (v,v') and (v,v'') , and $E_L(v',v'')$. A pair of edges (v,v') and (v,v'') are called the Δ -edges of v . We say that v is connected to C_i in terms of $\Delta(v)$.

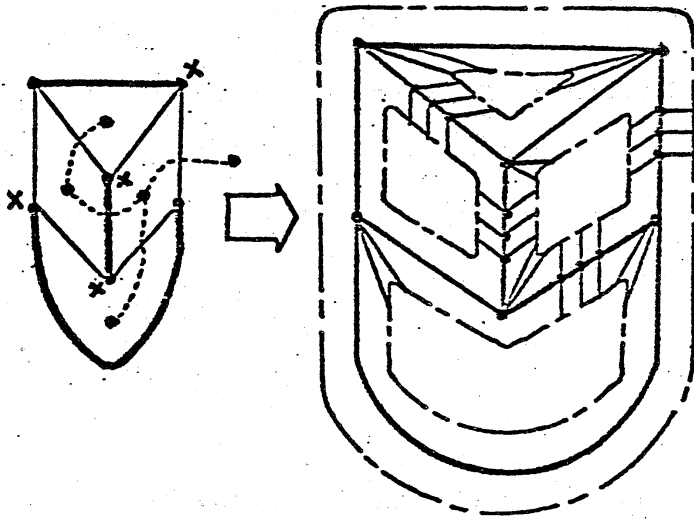


Figure 1. An example of the graph transformation from G to G_2 (In the figure $\bullet\text{---}\bullet$ and --- , respectively, denote an element of M and an inner ring, and both $\bullet\text{---}\bullet$ and $\bullet\text{---}\text{---}\bullet$ edges).

Lemma 2.

G has a node cover of size less than or equal to k if and only if there exists a subset $X \subset V_2$ with $|X| \leq k+r-1$ such that G_2-X is an s-p graph.

Proof. Let $N \subset V$ be a node cover for G with $|N| \leq k$. Define $X_e \subset V_1$ for each $e=(u,v) \in E$ as follows:

$$X_e = \begin{cases} \phi & \text{when } e \notin E_T, \\ \{w\} & \text{when } e \in E_T \text{ and } |\{u,v\} \cap N| = 1, \\ \{v_2\} & \text{when } e \in E_T \text{ and } \{u,v\} \subset N, \end{cases}$$

where if $\{u,v\} \cap N = \{u\}$ ($\{v\}$), then $w=v_3$ ($w=v_1$) for

$V_P(u,v) = \{v_0=u, v_1, v_2, v_3, v_4=v\}$. Let

$$X_1 = \bigcup_{e \in E} X_e \quad (= \bigcup_{e \in E_T} X_e)$$

and

$$X = N \cup X_1.$$

Note that $N \cap X_1 = \emptyset$. Then we have $|X| \leq k+r-1$, since $|E_T| = r-1$. It is easy to see that $G_2 - X$ is an s-p graph (see Figure 2).

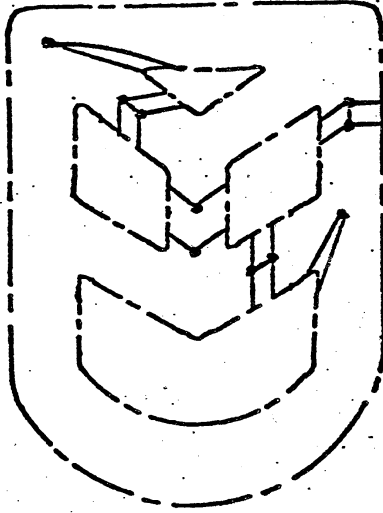


Figure 2. $G_2 - X$ for G_2 of Figure 1
(The node cover N that determines X consists of those nodes marked x in Figure 1).

Conversely, assume that an s-p graph can be formed by removing $(k+r-1)$ or fewer nodes of G_2 . Let $X \subset V_2$ be a minimal set satisfying that $G_2 - X$ is s-p and which maximizes, among all such sets, the number of nodes contained in V . By the construction of G_2 (in which any band of each inner ring is of width $(k+r+2)$), by Remark 1 and by the minimality of X , we can assume that

$$X \subset V_1.$$

Let

$$N = X \cap V$$

and let $P(u,v)$ be the T-path for arbitrary edge $e=(u,v) \in E_T$. Then

$$|X \cap \{v_1, v_2, v_3\}| \geq 1$$

for $V_{P_a}(u,v) = \{v_1, v_2, v_3\} \subset V_{P(u,v)}$, since otherwise $G_2 - X \supset K_4$.

Therefore, since we have $|E_T| = r-1$,

$$|N| \leq k.$$

Let $G_3 = (V_3, E_3)$ denote $G_2 - X$ for simplicity.

Lemma 10.

G_3 satisfies the following (1) and (2).

$$(1) \quad \delta_{G_3}(v) = 2 \text{ for } \forall v \in V \cap V_3.$$

(2) C_i and C_j are connected only in terms of exactly two c -paths (including the case when there exists an edge connecting these two c -paths) for any pair of distinct inner rings C_i and C_j that are T -connected in G_2 .

By Lemma 10, we can prove that N is a node cover for G .

Q.E.D.

Lemma 2 and the fact that G_2 can be constructed in time polynomial in the size of input establish the main theorem.

Theorem 2.

P_{PSPNR} is NP-complete.

We will prove Lemma 10 by way of Lemma 3 through Lemma 9 in the rest of this section.

We have $\delta_{G_3}(v) \geq 2$ for $\forall v \in V_3$. And if $v \in V \cap V_3$ then there exist $\Delta(v)$ and the inner ring to which v is connected in terms of $\Delta(v)$.

By the construction of G_2 , any two distinct inner rings C_i and C_j have no node in common. Thus G_3 has a path of length not less than 2 connecting C_i and C_j since $X \subset V_1$ and since G_3 is connected. Two blocks B and B' having a cutpoint u in common are said to be adjacent to each other at u .

Lemma 3.

Any block B of G_3 has at most two other blocks each of which is adjacent to B .

Let

$$K_\beta = (V_{K_\beta}, E_{K_\beta})$$

denote the graph shown in Figure 3, where

$$V_{K_\beta} = \{u, v, w_1, w_2\}, \text{ and}$$

$$E_{K_\beta} = \{(u, w_i), (v, w_i) : i=1, 2\} \cup \{(w_1, w_2)\}.$$

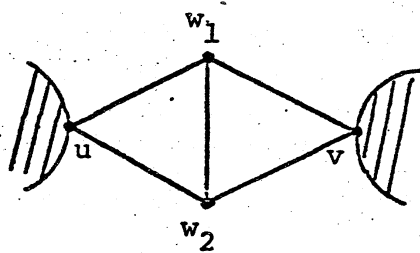


Figure 3. K_β .

Observe that a graph H is not s - p if K_β is a subgraph of H having both u and v as cutpoints of H (any reduction can not be applied to K_β since all of u, v, w_1 and w_2 are of node degree ≥ 3). Let

$$\tilde{K}_\beta = (V_{\tilde{K}_\beta}, E_{\tilde{K}_\beta})$$

denote the graph obtained by replacing each edge $(w, w') \in E_{K_\beta}$ with a disjoint path of length not less than 1 (that is, \tilde{K}_β is a graph homeomorphic to K_β). We say that a graph H has a pair of K_β -cutpoints if and only if H has a subgraph isomorphic to \tilde{K}_β having both u and v as cutpoints of H .

Lemma 4.

G_3 has no block which contains a pair of K_β -cutpoints.

Lemma 5.

Assume that G_3 has a triangle $\Delta(u)$ ($u \in V$). If $\delta_{G_3}(u)$

≥ 3 then u is a cutpoint of G_3 .

Lemma 6.

Assume that G_3 has a triangle $\Delta(u)$ ($u \in V$). Let C_i be the inner ring to which u is connected in terms of $\Delta(u)$, $\delta_{G_3}(u) \geq 3$ and $B(C_i)$ be the block of G_3 containing C_i . Then $B(C_i)$ has no other cutpoint of G_3 except u .

Let $G_{A_1} = (V_{A_1}, E_{A_1})$ ($V_{A_1} = A_1$) denote the subgraph of G_1 induced by A_1 , and let $\overline{E}_1 = E_1 - E_{A_1}$. Then let

$$IR(G_2) = (V_{IR(G_2)}, E_{IR(G_2)}),$$

called the inner ring tree, denote the subgraph of G_2 formed by deleting \overline{E}_1 from G_2 , where $E_{IR(G_2)} = E_2 - \overline{E}_1$. Let C and C' be a pair of inner rings that are T -connected in G_2 . We say that C and C' are π -connected, $C\pi C'$ for short, if and only if they are connected only in terms of exactly two c -paths (including the case when there exists an edge connecting these two c -paths). For any two distinct inner rings C and C' of a block, a sequence of distinct inner rings $C_{i_0} = C, C_{i_1}, \dots, C_{i_t} = C'$ ($t \geq 1$) of this block is called the π -connected sequence of C and C' if and only if $C_{i_j} \pi C_{i_{j+1}}$ for $j=0, \dots, t-1$.

Let D denote the set of all cutpoints of G_3 . Then, by the construction of G_2 and by the definition of X , we have

$$D \subset V_1 = V \oplus A_1.$$

Let us denote

$$D = D_V \oplus D_{A_1},$$

where $D_V = D \cap V$ and $D_{A_1} = D \cap A_1$. Then, by Lemmas 3, 5 and 6,

$$|D_V| \leq 2.$$

Let

$$X^{(1)} = X \cup D_V \quad \text{and} \quad G_3^{(1)} = G_2 - X^{(1)} \quad (= G_3 - D_V).$$

Note that $G_3^{(1)}$ is not always connected. Let B_1, B_2, \dots denote the blocks of $G_3^{(1)}$, and let

$$\mathcal{B}^{(1)} = \{B_1, B_2, \dots\}.$$

Hereafter, let us denote

$$G_3^{(i)} = (V_3^{(i)}, E_3^{(i)}) \quad \text{for } i=1, 2 \text{ and } 3.$$

Lemma 7.

$G_3^{(1)}$ satisfies the following (1), (2) and (3).

$$(1) \quad \delta_{G_3^{(1)}}(v) = 2 \quad \text{for } \forall v \in V \cap V_3^{(1)}.$$

$$(2) \quad D_{A_1} \subset D^{(1)} \subset A_1, \quad \text{where } D^{(1)} \text{ is the set of all cutpoints of } G_3^{(1)}.$$

$$(3) \quad \delta_{G_3^{(1)}}(u) \geq 2 \quad \text{for } \forall u \in V_3^{(1)}.$$

Lemma 8.

For any pair of distinct inner rings C and C' of any block $B_i \in \mathcal{B}^{(1)}$, there exists the π -connected sequence of C and C' .

Let

$$D^{(1)} = \{a_1, \dots, a_\theta\} \quad (\theta = |D^{(1)}|),$$

where the order of elements of $D^{(1)}$, denoted by the subscript, can arbitrarily fixed. Let $P(u, v)$ denote the T-path containing $a_1 \in D^{(1)}$.

Then, clearly,

$$a_1 \in V_{P_a}(u, v) = \{v_1, v_2, v_3\} \subset V_{P(u, v)} = \{u=v_0, v_1, v_2, v_3, v=v_4\}.$$

We define $X_{a_1} \subset V_1$ for the following two cases. Note that, for each $a_i \in D^{(1)}$, there exists exactly one T-path containing a_i and that if $a_i \neq a_j$ ($a_i, a_j \in D^{(1)}$) then they are contained in distinct T-paths.

(i) The case when $a_1=v_1$ (respectively when $a_1=v_3$).

By the definition of $G_3^{(1)}$,

$$\{v_0, v_2, v_3\} \subset X^{(1)} \quad (\{v_1, v_2, v_4\} \subset X^{(1)}).$$

Then let

$$X_{a_1} = X^{(1)} - \{v_2\}.$$

(ii) The case when $a_1=v_2$.

By the definition of $G_3^{(1)}$,

$$\{v_1, v_3\} \subset X^{(1)}.$$

Note that if $v_0 \notin X^{(1)}$ (if $v_4 \notin X^{(1)}$) then $\delta_{G_3^{(1)}}(v_0)=2$

($\delta_{G_3^{(1)}}(v_4)=2$) (that is, v_0 (v_4) is connected to

some inner ring only in terms of $\Delta(v_0)$ ($\Delta(v_4)$).

Then let

$$X_{a_1} = \begin{cases} X^{(1)} - \{v_1\} & \text{when } v_0 \in X^{(1)} \text{ and } v_4 \notin X^{(1)}, \\ X^{(1)} - \{v_3\} & \text{when } v_0 \notin X^{(1)} \text{ and } v_4 \in X^{(1)}, \\ (X^{(1)} - \{v_1\}) \cup \{v_0\} & \text{when } \{v_0, v_4\} \cap X^{(1)} = \emptyset, \\ X^{(1)} - \{v_1\} & \text{when } \{v_0, v_4\} \subset X^{(1)}. \end{cases}$$

Set

$$X_{a_0} = X^{(1)}$$

and, for each $j=2, \dots, \theta$, define recursively

$$X_{a_j} \subset v_1,$$

similarly to the definition of X_{a_1} , by replacing $X_{a_{j-2}}$ and $X_{a_{j-1}}$

with $X_{a_{j-1}}$ and X_{a_j} , respectively. Clearly, for each $j=1, \dots, \theta$,

we have either

$$|X_{a_{j-1}}| > |X_{a_j}|$$

or

$$|X_{a_{j-1}} \cap v| < |X_{a_j} \cap v| \text{ with } |X_{a_{j-1}}| = |X_{a_j}|.$$

Set

$$X^{(2)} = X_{a_\theta} \quad \text{and} \quad G_3^{(2)} = G_2 - X^{(2)}.$$

Then by the definition of $X^{(2)}$ and by Lemma 8 we have the next lemma.

Lemma 9.

$G_3^{(2)}$ satisfies the following (1), (2) and (3).

- (1) Any maximal connected component of $G_3^{(2)}$ is a block of $G_3^{(2)}$.
- (2) For any pair of distinct inner rings C and C' of any block of $G_3^{(2)}$, there exists the π -connected sequence of C and C' .
- (3) $\delta_{G_3^{(2)}}(v) = 2$ for $\forall v \in V \cap V_3^{(2)}$.

Lemma 9 shows that $G_3^{(2)}$ is a subgraph of the inner ring tree $IR(G_2)$.

For $w \in D_V$, let $B(w)$ denote the block of G_3 containing $\Delta(w)$, B be a block distinct from $B(w)$ and adjacent to $B(w)$ having w as a cutpoint. Let $(w, w^{(i_1)})$ and $(w, w^{(i_2)})$ be the two Δ -edges. Then we can show that any edge incident upon w and distinct from these two Δ -edges is contained in B . Therefore, $G_3^{(1)}$ consists of $|D_V|$ maximal connected components $B^{(1)}(w) = B(w) - \{w\}$ for $w \in D_V$ and another one C_Z which is determined by contracting $B(w)$ to w for $\forall w \in D_V$ in G_3 . Hence, by Lemma 9, $G_3^{(2)}$ consists of $(|D_V| + 1)$ blocks (or equivalently the same number of maximal connected components). That is, each $B^{(1)}(w)$ of $G_3^{(1)}$ is itself a block $B^{(2)}(w)$ of $G_3^{(2)}$ by Lemmas 6, 7 and 8, and C_Z of $G_3^{(1)}$ is changed into another block $B^{(2)}$ of $G_3^{(2)}$ in the process constructing $G_3^{(2)}$.

Now let us assume that $|D_V| \geq 1$ (In the case when $D_V = \emptyset$, we omit the procedure to construct $X^{(3)}$ described in the following, and set $X^{(3)} = X^{(2)}$ and $G_3^{(3)} = G_3^{(2)}$). Let

$$\mathcal{B}^{(2)} = \{B^{(2)}(w) : w \in D_V\} \cup \{B^{(2)}\} = \{B_1, \dots, B_{|D_V|+1}\}.$$

Let $B_i \in \mathcal{B}^{(2)}$ be arbitrarily chosen. By the definition of T^* or of $IR(G_2)$, there exist some $B_j \in \mathcal{B}^{(2)} - \{B_i\}$, some inner ring C_i of B_i and some inner ring C_j of B_j such that C_i and C_j are T -connected in G_3 . Then we say that B_j is a T -connected block of B_i and vice versa. Let $e=(u,v) \in E_T$ be the c -edge of C_i and C_j , and $P(u,v)$ be the corresponding T -path for e . Then

$$V_{P_a}(u,v) = \{v_1, v_2, v_3\} \subset X^{(2)},$$

where we set $u=v_0$ and $v=v_4$. Note that the following discussion does not depend on whether or not $D_V \cap \{u,v\}$ is empty. Define $X_{ij} \subset V_1$ as follows:

$$X_{ij} = \begin{cases} (X^{(2)} - \{v_1, v_2\}) \cup \{u\} & \text{when } \{u,v\} \cap X^{(2)} = \emptyset, \\ X^{(2)} - \{v_1, v_2\} & \text{when } \{u,v\} \cap X^{(2)} = \{u\}, \\ X^{(2)} - \{v_2, v_3\} & \text{when } \{u,v\} \cap X^{(2)} = \{v\}, \text{ and} \\ X^{(2)} - \{v_1, v_2\} & \text{when } \{u,v\} \subset X^{(2)}. \end{cases}$$

In the case when $|D_V| = 1$ define X_{ij} for $w \in D_V$ by setting $B_i = B^{(2)}(w)$.

In the case when $|D_V| = 2$, first define X_{ij} for $w \in D_V$ by setting $B_i = B^{(2)}(w)$. Set $B_s = B^{(2)}(w')$ for $w' \in D_V - \{w\}$ and let $B_t \in \mathcal{B}^{(2)}$ be a T -connected block of B_s . Then define X_{st} , similarly to the definition of X_{ij} , by replacing X_{ij} and $X^{(2)}$ with X_{st} and X_{ij} , respectively. Now we set

$$X^{(3)} = \begin{cases} X_{ij} & \text{when } |D_V| = 1, \text{ and} \\ X_{st} & \text{when } |D_V| = 2, \end{cases}$$

and clarify the relation between X and $X^{(3)}$. First,

$$|X_{ij}| \leq |X^{(2)}| - 1 \quad \text{and} \quad |X_{st}| \leq |X_{ij}| - 1,$$

so that

$$|X^{(3)}| \leq |X^{(2)}| - |D_V|,$$

where $|D_V| \leq 2$. Secondly,

$$|X| = |X^{(1)}| - |D_V| \text{ and } |X^{(1)}| \geq |X^{(2)}| ,$$

so that

$$|X| \geq |X^{(3)}| .$$

We have

$$|X| = |X^{(3)}|$$

only if we have the following (i) and (ii):

$$(i) \quad |X^{(1)}| = |X^{(2)}| .$$

(ii) $\{u, v\} \cap X^{(2)} = \emptyset$ and $\{u, v\} \cap X_{ij} = \emptyset$, respectively, for the c-edges $e=(u, v)$ in the definition of X_{ij} and of X_{st} .

If

$$|X^{(1)}| = |X^{(2)}|$$

then

$$|X^{(1)} \cap V| < |X^{(2)} \cap V|$$

as mentioned earlier. And (ii) implies that we have

$$|X_{ij} \cap V| = |X^{(2)} \cap V| + 1 \text{ and } |X_{st} \cap V| = |X_{ij} \cap V| + 1,$$

so that

$$|X^{(3)} \cap V| = |X^{(2)} \cap V| + |D_V| .$$

By the definition of $X^{(1)}$, we have

$$|X^{(1)} \cap V| = |X \cap V| + |D_V| .$$

Therefore, if

$$|X| = |X^{(3)}|$$

then

$$|X \cap V| < |X^{(3)} \cap V| .$$

Let

$$G_3^{(3)} = G_2 - X^{(3)} .$$

Then, clearly, we have the following statements (a) and (b).

- (a) $G_3^{(3)}$ is a nonseparable subgraph of the inner ring tree $IR(G_2)$ and satisfies the following (i) and (ii):

$$(i) \quad \delta_{G_3^{(3)}}(v) = 2 \text{ for } \forall v \in V_3^{(3)} \cap V.$$

(ii) For any pair of distinct inner rings C and C' of G_2 , there exists the π -connected sequence of C and C' in $G_3^{(3)}$.

$$(b) \quad |X| \geq |X^{(3)}|, \text{ where if } |X| = |X^{(3)}| \text{ then } |X \cap V| < |X^{(3)} \cap V|.$$

$G_3^{(3)}$ is essentially the same form as the G_2 - X constructed in the first half of the proof of Lemma 2, and it is easy to see that $G_3^{(3)}$ is s-p. Now the discussion so far is summarized as follows.

The proof of Lemma 10. Assume that either (1) or (2) of the lemma is false. Then we can define $X^{(3)} \subset V_1$ which is distinct from X and which satisfies the above statements (a) and (b), contradicting our choice of X . Q.E.D.

4. Concluding remark.

The problems P_{PSPNR} and P_{SPNR} discussed in this paper and in [15], respectively, are not included in the class of problems with properties that are hereditary on induced subgraphs, while a large number of node-deletion NP-complete problems with hereditary properties have been presented by M. S. Krishnamoorthy and N. Deo [11], and by M. Yannakakis [17, 18].

Acknowledgement.

The first author was supported by the Grant in Aid for Scientific Research of the Ministry of Education, Science and Culture of Japan under Grant: Encouraged Research (A) 475284.

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